MySQL PERFORMANCE CODING: FROM SOUP TO NUTS

By Jay Pipes Community Relations Manager, North America MySQL, Inc.



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Who the heck am I?

THE EXPERT'S VOICE Michael Kruckenberg and Jay Pipes Apress[®]

- Just some dude who works at MySQL
- Working with PHP and MySQL for 6+ years
- Oh, I wrote a book, too…
- Other than that, I'm semi-normal, with wife, two cats, two dogs, blah, blah, blah



A quick survey...

- 3.23? 4.0? 4.1? 5.0? 5.1?
- MyISAM? InnoDB? Archive? Memory?
- Replication? Cluster?
- PHP 4? PHP 5? PHP 6?
- libmysql? Native Driver for MySQL?
- ext/mysql? ext/mysqli? PDO?
- Oracle? PostgreSQL? DB2? MSSQL? SQLite?



The big three topics

The Schema



The Code



The Server





But first, some general stuff...

- Performance != Scalability
- Benchmarking
- Overview of MySQL System Architecture
- Overview of the MySQL Query Cache
- The Scan vs. Seek Choice



Performance !== Scalability

- Typically, we speak of performance when we talk about *response times* for a web page, an SQL statement, etc
- Scalability comes up when we talk about *throughput*, or the number of concurrent requests a node can serve within a certain timeframe
 or the size of the *data* increases



Benchmarking

- Very important to benchmark both performance and scalability
 - So, test with multiple concurrency levels and varying dataset sizes
- Toolbox
 - mysqlslap
 - Apache Bench (ab)
 - sysbench
 - Custom
 - MyBench, Jmeter/Ant, Shell scripts...



MySQL system architecture





MySQL Performance Coding

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MySQL system architecture notes

- Highly coupled subsystems
- Emphasis on connection-based memory allocation (as opposed to global)
- Caching on many different levels
- Storage engine layer is both blessing and curse
- Optimizer is cost-based, simplistic, and must be guarded against
- Efforts to modularize ongoing



The MySQL query cache

- You *must* understand application read/write ratio
- Internal design is a compromise between CPU usage and read performance
- Bigger query cache != better performance, even for heavy read applications
- Not a silver bullet!



The scan vs. seek choice

- A seek operation, generally speaking, jumps to a random place -- either on disk or in memory -- to fetch the data needed.
 - Repeat for each piece of data needed from disk or memory
- A scan operation, on the other hand, will jump to the start of a chunk of data, and sequentially read data -- either from disk or from memory -until the end of the chunk of data
 - For large amounts of data, scan operations tend to be more efficient than multiple seek operations



And finally...

- Learn to use EXPLAIN!
- Too big a topic for this session, so download my workbook and slides from OSCON
- http://jpipes.com/presentations/target-practice
 - /target-practice-workbook.pdf
 - /target-practice.pdf (or .odp)



MySQL Performance Coding

The Schema





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- Smaller, smaller, smaller!
 - Do you really need that BIGINT?
- More records in single page of memory, faster seeks & scans
- AUTO_INCREMENT is a good thing!
 - Generates a "hot spot" on disk and in memory
 - Look at 5.1.21 for InnoDB scaling patch
- Use appropriate data types



Appropriate data types

- INT UNSIGNED for IP addresses!
- Use VARCHAR carefully
 - Converted to CHAR when used in a temporary table
- Use TEXT sparingly
 - Consider separate tables
- Use BLOBs very sparingly
 - Use the filesystem for what it was intended



IPv4 addresses are INT UNSIGNED

CREATE TABLE Sessions (session_id INT UNSIGNED NOT NULL AUTO_INCREMENT ip_address_INT_UNSIGNED_NOT_NULL // Compared to CHAR(15)

- , ip_address INT UNSIGNED NOT NULL // Compared to CHAR(15)!!
- , session_data TEXT NOT NULL
- , **PRIMARY KEY** (session_id)
- , INDEX (ip_address)
-) ENGINE=InnoDB;



The INET_ATON() function reduces the string to a constant INT and a highly optimized range operation will be performed for:

SELECT * FROM Sessions
WHERE ip_address BETWEEN 3232235521 AND 3232235775



Normalize first, denormalize later. But...



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Partitioning

- Vertical partitioning
 - Splits tables with many columns into multiple tables
- Horizontal partitioning
 - Splits table with many rows into multiple tables
- Both are important for different reasons
- Partitioning in MySQL 5.1 is *horizontal partitioning*



Vertical partioning example

```
CREATE TABLE Users (
  user id INT NOT NULL AUTO INCREMENT
, email VARCHAR(80) NOT NULL
 display name VARCHAR(50) NOT NULL
 password CHAR(41) NOT NULL
 first name VARCHAR(25) NOT NULL
 last name VARCHAR(25) NOT NULL
  address VARCHAR(80) NOT NULL
 city VARCHAR(30) NOT NULL
  province CHAR(2) NOT NULL
  postcode CHAR(7) NOT NULL
  interests TEXT NULL
  bio TEXT NULL
  signature TEXT NULL
 skills TEXT NULL
  PRIMARY KEY (user id)
  UNIQUE INDEX (email)
  ENGINE=InnoDB;
```

```
CREATE TABLE Users (
  user id INT NOT NULL AUTO INCREMENT
, email VARCHAR(80) NOT NULL
 display name VARCHAR(50) NOT NULL
, password CHAR(41) NOT NULL
 PRIMARY KEY (user id)
 UNIQUE INDEX (email)
) ENGINE=InnoDB;
CREATE TABLE UserExtra (
  user id INT NOT NULL
, first name VARCHAR(25) NOT NULL
 last name VARCHAR(25) NOT NULL
 address VARCHAR(80) NOT NULL
, city VARCHAR(30) NOT NULL
, province CHAR(2) NOT NULL
 postcode CHAR(7) NOT NULL
 interests TEXT NULL
 bio TEXT NULL
 signature TEXT NULL
 skills TEXT NULL
 PRIMARY KEY (user id)
 ENGINE=InnoDB;
```



When vertical partitioning makes sense

- "Extra" columns are mostly NULL
- "Extra" columns are infrequently accessed
- When space in buffer pool is at a premium
 - Splitting the table allows main records to consume the buffer pages without the extra data taking up space in memory
 - Many more "main" records can fit into a single 16K InnoDB data page
- Need FULLTEXT on your text columns?



Vertical partioning example #2



```
CREATE TABLE Products (
  product id INT NOT NULL
, name VARCHAR(80) NOT NULL
  unit cost DECIMAL(7,2) NOT NULL
 description TEXT NULL
 image path TEXT NULL
, PRIMARY KEY (product id)
 INDEX (name(20))
) ENGINE=InnoDB; // Or MyISAM
CREATE TABLE ProductCounts (
  product id INT NOT NULL
, num views INT UNSIGNED NOT NULL
  num in stock INT UNSIGNED NOT NULL
 num on order INT UNSIGNED NOT NULL
 PRIMARY KEY (product id)
) ENGINE=InnoDB:
CREATE TABLE TableCounts (
  total products INT UNSIGNED NOT NULL
```

) ENGINE=MEMORY;



Vertical partitioning solves more problems

- Mixing *static* attributes with *frequently updated* fields in a single table?
 - Thrashing occurs with query cache. Each time an update occurs on any record in the table, all queries referencing the table are invalidated in the Query Cache
- Doing COUNT(*) with no WHERE on an indexed field on an InnoDB table?
 - Complications with versioning make full table counts very slow



Horizontal partitioning options

- Eli covered this well yesterday
 - (Thanks so much, Eli, I deleted four slides from my talk about it.)
- MySQL 5.1 Partitioning
 - (Will be) good for taking advantage of multiple disks
- Custom partitioning (sharding)
 Currently, the way to go



Indexes

- Speed up SELECTs, but slow down modifications
- Ensure indexes on columns used in WHERE, ON, GROUP BY clauses
- Always ensure JOIN conditions are indexed (and have identical data types)
- Be careful of the column order!
- Look for covering indexes



What makes a column ideal for indexing?

- Selectivity
 - % of distinct values in a column
 - S=d/n
 - Unique/primary always be 1.0
- If column has a low selectivity, it can still be put in a multi-column index
 - But, which part? Prefix? Suffix?



Remove redundant or poor indexes

SELECT
t.TABLE_SCHEMA AS `db`
, t.TABLE NAME AS `table`
, s.INDEX_NAME AS `inde name`
, s.COLUMN NAME AS `field name`
, s.SEQ_IN_INDEX `seq in index`
, s2.max columns AS `# cols`
, s.CARDINALITY AS `card`
, t.TABLE ROWS AS `est rows`
, ROUND(((s.CARDINALITY / IFNULL(t.TABLE_ROWS, 0.01)) * 100), 2) AS `sel %`
FROM INFORMATION_SCHEMA.STATISTICS s
INNER JOIN INFORMATION SCHEMA.TABLES t
ON s.TABLE SCHEMA = t.TABLE SCHEMA
AND s.TABLE_NAME = t.TABLE_NAME
INNER JOIN (
SELECT TABLE SCHEMA, TABLE NAME, INDEX NAME, MAX(SEQ IN INDEX) AS max columns
FROM INFORMATION SCHEMA.STATISTICS
WHERE TABLE_SCHEMA != 'mysql'
GROUP BY TABLE SCHEMA, TABLE NAME, INDEX NAME
) AS s2
ON s.TABLE SCHEMA = s2.TABLE SCHEMA
AND s.TABLE_NAME = s2.TABLE_NAME
AND S.INDEX NAME = S2.INDEX NAME
WHERE t.TABLE_SCHEMA != 'mysql' /* Filter out the mysql system DB */ AND t.TABLE_ROWS > 10 /* Only tables with some rows */
AND s.CARDINALITY IS NOT NULL /* Need at least one non-NULL value in the field */
AND (s.CARDINALITY / IFNULL(t.TABLE_ROWS, 0.01)) < 1.00 /* unique indexes are perfect anyway */
ORDER BY `sel %`, s.TABLE_SCHEMA, s.TABLE_NAME /* DESC for best non-unique indexes */
LIMIT 10;
http://forgo.my/ogl.com/opinp.cts/viouv.php?id=95

http://forge.mysql.com/snippets/view.php?id=85



Whoah. Some crap-ass indexes, huh?

TABLE_SCHEMA	TABLE_NAME	INDEX_NAME	COLUMN_NAME	SEQ_IN_INDEX	COLS_IN_INDEX	CARD	ROWS	SEL %
worklog	amendments	text	text	1	1	1	33794	0.00
planetmysql	entries	categories	categories	1	3	1	j 4171	0.02
planetmysql	entries	categories	title	2	3	1	j 4171	0.02
planetmysql	entries	categories	content	3	3	1	i 4171	i 0.02
sakila	inventory	idx store id film id	store id	i 1 i	2	1	4673	0.02
sakila	rental	idx fk staff id	staff_id	1	1	3	16291	0.02
worklog	tasks	title	title	1	2	1	3567	0.03
worklog	tasks	title	description	2	2	1	3567	0.03
sakila	payment	idx fk staff id	staff id	i <u>ī</u> i	1	6	15422	0.04
mysqlforge	mw_recentchanges	rc_ip	rc_ip		1	2	996	0.20



Effect of index column order

-> FROM Tag -> GROUP B	g2Projeci / project	:	DUNT(*) as num_tag	s
table	type	key	Extra	
	•		Using index; <mark>Usi</mark>	ng temporary; Using filesort
-> FROM Tag -> GROUP B	g2Project / tag;	:	(*) as num_project	s The Tag2Project Table:
table	type	key	Extra	CREATE TABLE Tag2Project (
Tag2Project +	index + INDEX pro	PRIMARY	<pre>tag INT UNSIGNED NOT NULL , project INT UNSIGNED NOT NULL , PRIMARY KEY (tag, project)) ENGINE=MyISAM;</pre>	
Records: 701 [How do you tell if this is happening?
-> FROM Tag -> GROUP B	g2Project / project	:	DUNT(*) as num_tag	<pre>Look for increases in: Created_tmp_tables and Created_tmp_disk_tables</pre>
table	type	key	Extra	status counter variables
Tag2Project	index	project	Using index	



Covering indexes

- When all columns needed from a single table for a SELECT are available in the index
- No need to grab the rest of the columns from the data (file or page)
 "Bookmark lookup" operation
- Important to know the data to index organization of the storage engine!



Non-clustered organization (MyISAM)



Clustered organization (InnoDB)



So, bottom line:

When looking up a record by a primary key, for a clustered layout/organization, the **lookup operation** (following the pointer from the leaf node to the data file) involved in a non-clustered layout **is not needed.**



Clustered indexes

- Very important to have as small a clustering key (primary key) as possible
 - Why? Because every secondary index built on the table will have the primary key appended to each index record
- If you don't pick a primary key (bad idea!), one will be created for you
 - And, you have no control over the key (this is a 6 byte number in InnoDB...)



MySQL Performance Coding

The Code





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Be a join-fu master!



Correlated subqueries shall die!!

- Don't think in terms of iterators, for loops, while loops, etc
- Instead, think in terms of sets
- Break complex SQL statements (or business requests) into smaller, manageable chunks



Set-wise thinking

"Show the last payment information for each customer"

- Many programmers think:
 - OK, for each customer, find the maximum date the payment was made get that payment record(s) (bad!)
- Think instead:
 - OK, I have 2 sets of data here. One set of last payments dates and another set containing payment information (so, how do I *join* these sets?)



The difference in execution plans?

-> (SELECT -> FROM p	/ment p .payment_date Γ MAX(payment_ payment customer_id=p	date)	id);					
select_type	table	typ	e possible_keys	key	ref	++- rows	Extra	-+
PRIMARY DEPENDENT SUE	p BQUERY payme	ALL nt ref	NULL idx_fk_customer_id,payment_date	NULL e payment_date			Using where Using index	
	customer_id,	MAX(paymo	ent_date) as last_order					
-> p.* -> FROM (-> SELECT -> FROM pa -> GROUP E ->) AS las -> INNER JO -> ON p.cus -> AND p.pa	customer_id, ayment 3Y customer_id st_orders DIN payment p stomer_id = la ayment_date =	st_order	<pre></pre>				++	
-> p.* -> FROM (-> SELECT -> FROM pa -> GROUP E ->) AS las -> INNER JC -> ON p.cus -> AND p.pa select_type	customer_id, ayment 3Y customer_id st_orders DIN payment p stomer_id = la ayment_date = ++ table	st_orders last_orde type	<pre>customer_id ers.last_order; possible_keys </pre>	key	ref		++ rows ++	
-> p.* -> FROM (-> SELECT -> FROM pa -> GROUP E ->) AS las -> INNER JC -> ON p.cus -> AND p.pa	customer_id, ayment 3Y customer_id st_orders DIN payment p stomer_id = la ayment_date =	st_order	<pre>s.customer_id ers.last_order; possible_keys NULL idx_fk_customer_id,payment_date</pre>	key NULL payment_date idx fk customer i	NULL customer_id,	,last_orde	599	


The difference in performance

<pre>mysql> SELECT -> p.* -> FROM payment p -> WHERE p.payment_date = -> (SELECT MAX(payment_date) -> FROM payment -> WHERE customer_id=p.customer_id);</pre>							
· -	customer_id	. – .		amount	payment_date	last_update	1
<snip> 16049 +</snip>				2.99	2005-08-23 11:25:00 +	2006-02-15 19:24:13 +	 .+
-> FROM µ -> GROUP ->) AS li -> INNER -> ON p.co	T customer_id, payment BY customer_ic ast_orders JOIN payment p ustomer_id = la payment_date =	d ast_orders.c last_orders	ustomer_id	-	+	+	-+
payment_id	customer_id				payment_date	' last_update +	 -+
<pre></pre>	599	2	15725	2.99	2005-08-23 11:25:00	2006-02-15 19:24:13	1
623 rows in so	et (0.09 sec)	+			+	+	+



A word on too many joins

- Don't try to do joins on >8 tables
 - Especially with MySQL < 5.0</p>
 - Sometimes optimizer can try too hard to find optimal plan
- Use a "temp table reduction" recipe
 - Especially important for AND conditions on many-to-many relation tables
- Or for small, static lookups, use ENUM (or SET for many-to-many)



Operating on indexed column w/ function

<pre>mysql> EXPLAIN SELECT * FROM film WHERE title LIKE 'Tr%'\G ************************************</pre>				
<pre>id: 1 select_type: SIMPLE table: film type: range possible_keys: idx_title key: idx_title</pre>	Nice. In the top query, we have a fast range access on the indexed field			
key_len: ref: rows: Extra:	NULL			

mysql> EXPLAIN SELECT * FROM	1 fil	Lm WHERE	LEFT(titl	.e,2) =	'Tr' ∖G
*****	1,	···· ****	********	******	*****
	- I I	0w .			

id: 1
select_type: SIMPLE
table: film
type: ALL
possible_keys: NULL
key: NULL
key_len: NULL
ref: NULL
rows: 951
Extra: Using where

Oops. In the bottom query, we have a slower full table scan because of the function operating on the indexed field (the LEFT() function)



Operating on indexed column w/ function #2

<pre>SELECT * FROM Orders WHERE TO_DAYS(CURRENT_DATE()) - TO_DAYS(order_created) <= 7;</pre>	Not a good idea! Lots o' problems with this		
<pre>SELECT * FROM Orders WHERE order_created >= CURRENT_DATE() - INTERVAL 7 DAY;</pre>	Better Now the index on order_created will be used at least. Still a problem, though		

SELECT order_id, order_created, customer			
FROM Orders	Best. Now the query cache can		
WHERE order_created	cache this query, and given no		
	updates, only run it once a day		

replace the CURRENT_DATE() function with a constant string in your programming language du jour... for instance, in PHP, we'd do:

```
$sql= "SELECT order_id, order_created, customer FROM Orders WHERE
order_created >= '" .
date('Y-m-d') . "' - INTERVAL 7 DAY";
```



Calculated fields example

CREATE TABLE Customers (
 customer_id INT NOT NULL
, email VARCHAR(80) NOT NULL
// more fields
, PRIMARY KEY (customer_id)
, INDEX (email(40))
) ENGINE=InnoDB;
// Bad idea, can't use index
// on email field
SELECT *
FROM Customers

WHERE email LIKE '%.com';

// So, we enable fast searching on a reversed field // value by inserting a calculated field **ALTER TABLE** Customers ADD COLUMN rv email VARCHAR(80) NOT NULL; // Now, we update the existing table values **UPDATE** Customers **SET** rv email = **REVERSE**(email); // Then, we create an index on the new field **CREATE INDEX** ix rv email **ON** Customers (rv email); // Then, we make a trigger to keep our data in sync **DELIMITER** ;; **CREATE TRIGGER** trg bi cust **BEFORE INSERT ON** Customers FOR EACH ROW BEGIN **SET NEW**.rv email = **REVERSE**(**NEW**.email); END ;; // same trigger for BEFORE UPDATE... // Then SELECT on the new field... WHERE rv email LIKE CONCAT(REVERSE('.com'), '%');



Using stored procedures

- Question: where does the stored procedure compile cache live?
- Don't use stored procedures for simple SELECTs
- Use for:
 - ETL or complex collections of SQL
 - Repeated execution of statement
 - Batch operations



UPDATEs and DELETEs

- Avoid DELETE, especially in MyISAM
 - Use a deleted_rows table
 - Insert rows into the table, then do batched DELETEs
- Have lots of UPDATEs?
 - Insert them into memcache bucket, then periodically update the tables...



MySQL Performance Coding

The Server





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SHOW STATUS and SHOW VARIABLES

SHOW STATUS

- Counter variables (lots of `em)
- Count reads, writes, threads, etc.
- SHOW VARIABLES
 - Your configuration variables
- Both take a LIKE clause, for example:





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Server variable guidelines

- Be aware of what is *global* vs *per thread*
- Make small changes, then test
- Often provide a quick solution, but temporary
- key_buffer_size != innodb_buffer_pool_size
- Memory is cheapest, fastest, easiest way to increase performance
 - But... bigger buffers aren't always a good thing!



Important settings

- **key_buffer_size** (global, MyISAM only)
 - Main MyISAM key cache (blocks of size 1K)
 - Watch for Key_blocks_unused approaching 0
- table_cache (global)
 - Number of simultaneously open file descriptors
 - < 5.1 contains meta data about tables and file descriptor
 - >= 5.1 Split into table_open_cache
- myisam_sort_buffer_size (global, MyISAM only)
 - Building indexes?
 - Set this as high as possible



Average table scan and key cache hit ratio

 Examine Handler_read_rnd_next/Handler_read_rnd for average size of table scans

۰;

mysql> SHOW STATUS LIKE	
Variable_name	Value
<pre>+ Handler_read_rnd Handler_read_rnd_next +</pre>	2188 217247

 Examine Key_read_requests/Key_reads for your MyISAM key cache hit ratio

mysql> SHOW STATUS L	
	Value
Key_read_requests Key_reads	10063 98 ++



Important settings for InnoDB

innodb buffer pool size

- Main InnoDB cache for both data and index pages (16K page)
- If you have InnoDB-only system, set to 60-80% of total memory
- Watch for Innodb buffer pool pages free approaching 0
- innodb_log_file_size
 - Size of the actual log file
 - Set to 40-50% of innodb buffer pool size
 - Bigger means longer recovery, but less disk I/O due to less checkpoint flush activity



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Important settings for InnoDB

• innodb_log_buffer_size

- Size of double-write log buffer
- Set < 16M (recommend 1M to 8M)</p>

innodb_flush_method

- Determines how InnoDB flushes data and logs
- defaults to fsync()
- If getting lots of Innodb_data_pending_fsyncs
 - Consider O_DIRECT (Linux only)
- Other ideas
- Get a battery-backed disk controller with a write-back cache
- Set innodb_flush_log_at_trx_commit=2 (Risky)



Important settings for InnoDB

 Examine Innodb_buffer_pool_reads vs Innodb_buffer_pool_read_requests for the cache hit ratio

mysql> SHOW STATUS LIKE 'Ir	nodb_buffe	<pre>er_pool_read%';</pre>
Variable_name		Value
<pre> Innodb_buffer_pool_read_r Innodb_buffer_pool_reads</pre>	requests 	5415365 34260
mysql> SHOW STATUS LIKE 'Qc%'	;	+
Qcache_free_blocks 1 Qcache_hits 6 Qcache_inserts 1		Examine C for the que Ensure QC is low
Qcache_lowmem_prunes 0	41 9 • 241 +	Ensure Qc • if not, F

- Examine Qcache_hits/Questions for the query cache hit ratio
- Ensure Qcache_lowmem_prunes is low
- Ensure Qcache_free_blocks = 1
 - if not, FLUSH QUERY CACHE



Yo, we're hiring.



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- 5 weeks vacation

